

Weighted Pólya counting, a.k.a. Pólya-Redfield enumeration

Notice that the example of

$G = \langle \tau = (1, 2, 3, 4) \rangle \curvearrowright$ (colorings of $[4]$)
is the same as w/ 2 colors (e.g. black+white)

$G \curvearrowright$ subsets of $[4]$

(Just ~~use~~ choose subset of black vertices.)

But we also know that an action on subsets like this preserves size of subset; e.g. we looked at

$G \curvearrowright$ size 2-subsets of $[4]$

which in the language of colorings would be same as

$G \curvearrowright$ colorings of $[4]$ w/ 2 vertices white + 2 vertices black.

In general we might want to keep track of precise number of each color used, as in:

Q: Up to rotation, how many colorings of vertices of square ^{have} exactly 2 red, 1 blue, 1 green vertex?

2/1 To answer this question, we need more notation.

for a coloring $f: X \rightarrow Y = \{1, 2, \dots, k\}$ define

monomial $\vec{y}^f := \prod_{x \in X} y_{f(x)} \in \mathbb{C}[y_1, y_2, \dots, y_k]$

e.g. ~~square~~ $R - R$.

Coloring $f = \begin{matrix} 1 & 1 \\ B & G \end{matrix} \rightsquigarrow y_1^2 y_2 y_3$ if we decide $\begin{cases} R=1 \\ B=2 \\ G=3 \end{cases}$

Notice: If $G \curvearrowright X$ then $\vec{y}_f = \vec{y}_{g \cdot f} \quad \forall g \in G$.

DEF'N Let $G \curvearrowright X$ and hence on colorings Y^X .
 The pattern inventory polynomial of $G \curvearrowright Y^X$ is

$$P(y_1, y_2, \dots, y_k) = \sum_{\Theta} \bar{y}^{\Theta} \in \mathbb{C}[y_1, \dots, y_k]$$

where the sum is over all orbits Θ of $G \curvearrowright Y^X$
 and $\bar{y}^{\Theta} := \bar{y}^f$ for any coloring $f \in \Theta$.

e.g. For $G = \langle \sigma = (1, 2, 3, 4) \rangle \curvearrowright X = [4]$ and $Y = \overset{\text{"}}{\Sigma} \{0, 1, 2\} = \{1, 2\}$,

$$P = 1y_1^4 + 1y_1^3y_2 + 2y_1^2y_2^2 + 1y_1y_2^3 + 1y_2^4$$

Given the pattern inventory poly. P , we can then answer questions like: "how many (symmetry classes of) colorings use 2 white + 2 black vertices?" by extracting coefficients.

So our goal will now be to give a formula for $P(y_1, \dots, y_n)$.
 To do that, we need to keep track of more refined cycle information of elements $g: X \rightarrow X, g \in G$.

Set $c_i(g) := \# \underset{\text{cycles of size } i}{\text{cycles of }} g: X \rightarrow X$.

DEF'N The cycle index polynomial of $G \curvearrowright X$ is

$$\mathcal{Z}_G(t_1, t_2, \dots, t_n) = \frac{1}{\#G} \sum_{g \in G} \prod_{i=1}^n t_i^{c_i(g)} \in \mathbb{C}[t_1, \dots, t_n]$$

This is the key to Poly-counting!

E.g. With $G = \langle \sigma = (1, 2, 3, 4) \rangle \cap X = [4]$, have

$$Z_G = \frac{1}{4} \left(\underbrace{t_1^4}_{\sigma = (1)(2)(3)(4)} + \underbrace{2t_4}_{\sigma = (1, 2, 3, 4)} + \underbrace{t_2^2}_{\sigma = (1, 3)(2, 4)} \right)$$

Thm (Pólya-Redfield enumeration theorem)

The pattern inventory polynomial of $G \cap X$ is

$$P = Z_G \left(\sum_{i \in Y} y_i, \underbrace{\sum_{i \in Y} y_i^2}_{t_1}, \underbrace{\sum_{i \in Y} y_i^3}_{t_2}, \dots, \underbrace{\sum_{i \in Y} y_i^n}_{t_n} \right).$$

E.g. Let $G = \langle \sigma = (1, 2, 3, 4) \rangle \cap X = [4]$ and consider set of colors $\gamma = \{R, G, B\} = \{1, 2, 3\}$.

$$\text{Then } P = \frac{1}{4} \left((y_1 + y_2 + y_3)^4 + 2(y_1^4 + y_2^4 + y_3^4) + (y_1^2 + y_2^2 + y_3^2)^2 \right)$$

$$\begin{aligned} &= \dots = y_1^4 + y_2^4 + y_3^4 + y_1^3 y_2 + y_1^3 y_3 + y_2^3 y_1 + y_2^3 y_3 \\ &\quad \stackrel{\text{lots of terms}}{+} y_3^3 y_2 + y_3^3 y_2 + 2(y_1^2 y_2^2 + 2y_1^2 y_3^2 \\ &\quad + 2y_2^2 y_3^2 + 3y_1^2 y_2 y_3 + 3y_2^2 y_1 y_3 + 3y_3^2 y_1 y_2). \end{aligned}$$

To figure out how many colorings have 2 R, 1 B, 1 G, we extract coeff. of $y_1^2 y_2 y_3$ from P;

$$\text{A: } [y_1^2 y_2 y_3] P = 3 \text{ colorings w/ 2R, 1B, 1G.}$$

Note: Setting $y_i = 1$ for all $i \in Y$, we recover the unweighted Pólya counting formula for total number of colorings (ignoring patterns).

2/9

Pf of Polya-Redfield Thm:

The proof is very similar to unweighted result; we just need to make sure we keep track of weights.

- First observe that for any orbit O of $G \curvearrowright Y^X$,

$$\vec{y}^O = \sum_{f \in O} \vec{y}^f / \#O_f, \text{ so}$$

$$P = \sum_O \vec{y}^O = \sum_{f \in Y^X} \frac{\vec{y}^f}{\#O_f} = \frac{1}{\#G} \sum_{f \in Y^X} \#G_f \cdot \vec{y}^f,$$

whereas before we used the Orbit-Stabilizer Thm.

By the same "summing over rows" vs.

"summing over columns" trick, applied to matrix

$$M(g, f) = \begin{cases} \vec{y}^f & \text{if } g \cdot f = f \\ 0 & \text{otherwise} \end{cases}, \text{ get that}$$

$$P = \sum_O \vec{y}^O = \frac{1}{\#G} \sum_{g \in G} \sum_{f \in (Y^X)^g} \vec{y}^f$$

- So we again need to think about $(Y^X)^g$ for $g \in G$.

Recall that $f \in (Y^X)^g \Leftrightarrow f(x) = f(x')$ whenever

x, x' belong to same cycle of $g: X \rightarrow X$

e.g. $g = (x_1, \underbrace{x_2, x_3}_{\text{color all red or all blue}}, x_4) \underbrace{(x_5, x_6)}_{\text{color red or blue}} : X \rightarrow X$

$$\sum_{f \in (Y^X)^g} \vec{y}^f = (y_1^3 + y_2^3 + y_3^3) \cdot (y_1 + y_2 + y_3) \cdot (y_1^2 + y_2^2 + y_3^2)$$

- So in general $\sum_{f \in (Y^X)^g} \vec{y}^f = \prod_{\substack{\text{cycles} \\ \text{of } g: X \rightarrow X}} \sum_{y \in Y} y_i^{\text{size of } c}$

This precisely means $P(y_1, \dots, y_K) = Z_G(\sum_i y_i, \sum_i y_i^2, \dots, \sum_i y_i^n)$.

Cor Let $G \sim X$. Then,

$$\sum_{k=0}^n \#(\text{orbits of } G \text{ w/ size } k \text{ subsets of } X) \cdot t^k = Z_G(1+t, 1+t^2, \dots, 1+t^n).$$

Pf: Use 2 colors in weighted polygon counting. \blacksquare

2/9 E.g. Recall that a ^(simple) graph consists of a vertex set V and a set of edges E , unordered pairs of vertices,

$$G = \begin{array}{c} 1 \\ \swarrow \quad \searrow \\ 2 \quad 3 \\ | \quad \diagdown \\ 4 \quad 5 \end{array} \Rightarrow V = [5], E = \{\{1,2\}, \{2,3\}, \{4,5\}\}$$

Q: How many graphs G with vertex set $V = [n]$?

A: $2^{\binom{n}{2}}$ since there are $\binom{n}{2}$ possible edges, and we can choose any subset of edges.

But... what if we want to count unlabeled graphs, i.e., graphs up to isomorphism?

DEFN An isomorphism between graphs $G = (V, E)$ and $G' = (V', E')$ is a bijection $\phi: V \rightarrow V'$ on vertices s.t. $\{i, j\} \in E \Leftrightarrow \{\phi(i), \phi(j)\} \in E'$.

$$\text{e.g. } \begin{array}{c} 1 \\ \swarrow \quad \searrow \\ 2 \quad 3 \end{array} = G \sim \begin{array}{c} 2 \\ \swarrow \quad \searrow \\ 1 \quad 3 \end{array} = G'$$

Q: How many isomorphism classes of graphs w/ n vertices are there?

Even better, what is $\sum_{G, \text{ graph on } n \text{ vertices}} t^{\#\text{edges}(G)}$?

A: By weighted Pólya counting, answer is

$$Z_G(1+t_1, 1+t^2, \dots, \cancel{t^{n-1}}, 1+t^{(2)})$$

where $G = S_n \curvearrowright X = \{\text{size 2 subsets of } [n]\}^3$.

\hookrightarrow WARNING! Not $X = [n]^3$

Let's first consider case $n=3$:

cycle type	$\sigma \in S_n$ w/ this type	cycle structure of $\sigma: X \rightarrow X$	monomial t^σ	# $\sigma \in S_n$ w/ type λ
(1, 1, 1)	$e = (1)(2)(3)$	(123)(132)(231)	t_1^3	1
(2, 1)	(1, 2) (3)	(123, 132, 231)	$t_2 t_1$	3
(3)	(1, 2, 3)	(123, 231, 312)	t_3	2

$$\text{So } Z_G(t_1, t_2, t_3) = \frac{1}{3!} (t_1^3 + 3t_2 t_1 + t_3)$$

$$\begin{aligned} \text{and } Z_G(1+t, 1+t^2, 1+t^3) &= \frac{1}{6} ((1+t)^3 + 3(1+t^2)(1+t) + 2(1+t^3)) \\ &= t^3 + t^2 + t + 1 \end{aligned}$$

$\Delta \quad \therefore \quad \hookrightarrow$ g.f. of graphs on 3 vertices, by #edges.

$n=4$: λ	$\sigma \in S_n$	cycle structure of $\sigma: X \rightarrow X$	t^σ	# $\sigma \in S_n$
(1, 1, 1, 1)	$e = (1)(2)(3)(4)$	(12)(13)(14)(23)(24)(34)	t_1^6	1
(2, 1, 1)	(1, 2)(3)(4)	(12)(13, 23)(14, 24)(34)	$t_2^2 t_1^2$	$\binom{4}{2} = 6$
(2, 2)	(1, 2)(3, 4)	(12)(13, 24)(14, 23)(34)	$t_2^2 t_1^2$	$\binom{4}{2}/2 = 3$
(3, 1)	(1, 2, 3)(4)	(12, 23, 13)(14, 24, 34)	t_3^2	$4 \cdot 2 = 8$
(4)	(1, 2, 3, 4)	(12, 23, 34, 14)(13, 24)	$t_4 t_2$	$3! = 6$

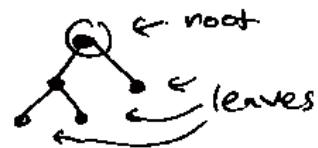
$$\text{So } Z_G(t_1, t_2, t_3, t_4) = \frac{1}{4!} (t_1^6 + 6t_2^2 t_1^2 + 8t_3^2 + 6t_4 t_2)$$

$$\begin{aligned} \text{and } Z_G(1+t, 1+t^2, 1+t^3, 1+t^4) &= \frac{1}{24} ((1+t)^6 + 9(1+t^2)^2(1+t)^2 \\ &\quad + 8(1+t^3)^2 + 6(1+t^4)(1+t^2), \\ &= t^6 + t^5 + 2t^4 + 3t^3 + 2t^2 + t + 1 \end{aligned}$$

$\Delta, \square, \square\square, \square\sqcup\sqcup, \square\sqcup\sqcup\sqcup, \square\sqcup\sqcup\sqcup\sqcup \quad \hookrightarrow$ g.f. for graphs on 4 vertices,

Cultural aside on trees: Polya developed Polya counting to enumerate trees, motivated by problems in molecular chemistry!

A rooted binary tree looks like:
each node has 2 or 0 children

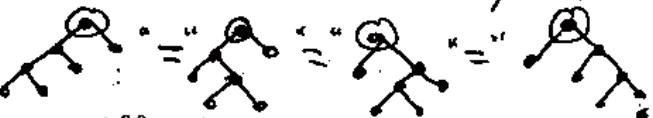


rooted binary trees w/ $n+1$ leaves = Catalan number. $C_n = \frac{1}{n+1} \binom{2n}{n}$ & sensat!

$$\text{e.g. } C_3 = 5$$



But what is we wanted to count "structurally different" binary trees, i.e.



Let $a_n := \#$ structurally different rooted binary trees w/ $n+1$ leaves

$$n = 0, 1, 2, 3, 4, 5, \dots$$

$$C_n = 1, 1, 2, 5, 14, 42, \dots$$

$$a_n = 1, 1, 1, 2, 3, 6, \dots$$

$$\text{Set } C(x) := \sum_{n \geq 0} C_n x^n \text{ and } A(x) := \sum_{n \geq 0} a_n x^n$$

$$\text{We saw } C(x) = 1 + x \cdot C(x)^2 \quad \text{account for symmetry!}$$

$$\text{Polya counting} \Rightarrow A(x) = 1 + \frac{x}{2} (A(x)^2 + A(x^2))$$

Even leads to asymptotics for all trees!

Thm (Offner, 1948) Let $t_n := \#$ unlabeled, unrooted trees on n vertices

$$\text{Then } t_n \sim C \alpha^n n^{-5/2} \text{ w/ } \alpha \approx 2.955 \dots$$

$$C \approx 0.5349 \dots$$

$$\text{Compare: } \frac{n^{n-2}}{n!} \sim \frac{1}{\sqrt{2\pi}} e^{-n} n^{-5/2} \approx 2.71 \dots$$

Cayley's formula for ~~labeled~~ [→] ~~labeled~~ trees!