

Fall 2019, UMN Math 8668:

Combinatorial Theory (1st sem. intro grad comb.)

Instructor: Sam Hopkins, shopkins@umn.edu

website: umn.edu/~shopkins/classes/8668.html

9/4/19

Class info:

- meets MWF 2:30-3:20, Ford Hall 360
- office hrs: MW 10-11am, Vincent Hall 204
(or by appointment if you email me...)
- Text:
R. Stanley's "Enumerative Combinatorics, Vol 1"
(link to pdf on website)
HW problems mainly from here
I'll also heavily use notes of F. Ardila (they are great! :))
(also linked to on website)
- Grading: There will be 3 HW's (roughly: Oct, Nov, Dec)
Beyond that, show up to class and be engaged
(e.g. ask questions) to get a good grade.
- This is the 1st half of a year-long sequence.
8669 will be taught by Chris Fraser in the Spring

What is this class about?

We want to count combinatorial objects

(subsets, multisets, partitions, graphs, etc.)

and more generally understand their structure
(e.g. partial order structures).

W11

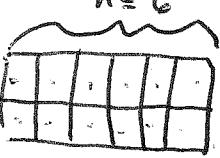
What is a "good" answer to a counting problem?

It depends! (For instance, on what we want to do w/ answer.)

Rather than try to formalize a notion of "good answer"
(let's explore what answers can look like in an example.)

Ex:

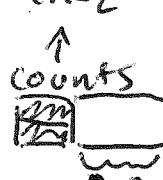
Let $a_n = \#$ tilings of $2 \times n$ rectangle by dominoes



either or

n	a_n	rectangle	tilings
0	1		
1	1		
2	2		,
3	3		, ,
4	5		, , , ,

① Recurrence! $a_n = a_{n-1} + a_{n-2}$ for $n \geq 2$, w/ $a_0, a_1 = 1$



This is the same recurrence as the Fibonacci numbers

$$F_n = F_{n-1} + F_{n+2} \text{ for } n \geq 2$$

Usually F_i are indexed w/ $F_0 = 0, F_1 = 1$

n	a_n	F_n
0	1	0
1	1	1
2	2	1
3	3	2
4	5	3

← So $a_n = F_{n+1}$,

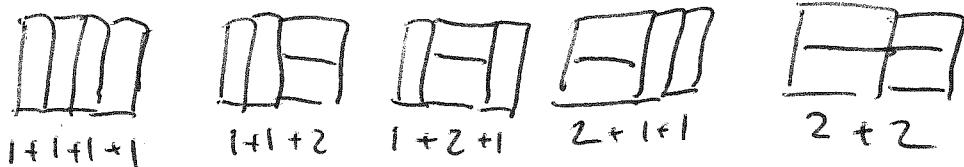
i.e., we're just considering Fibonacci numbers w/ slightly different indexing...)

This recurrence lets us compute a_n , but it might take a while + we don't get a sense of growth rate of a_n .

② Explicit formula as a summation

Observe $a_n = \#\{\text{Sequences of 1's and 2's summing to } n\}$

e.g. $n=4$



$$\Rightarrow a_n = \sum_{k=0}^{\lfloor n/2 \rfloor} \#\{\text{Sequences of } k \text{ 2's and } n-2k \text{ 1's}\}$$

$$\text{e.g. } 1 + \underbrace{2+1+1}_{\uparrow \uparrow \uparrow} + 2 + 1 + 2 = n = 10$$

$$= \sum_{k=0}^{\lfloor n/2 \rfloor} \binom{(n-2k)+k}{k} = \sum_{k=0}^{\lfloor n/2 \rfloor} \binom{n-k}{k}$$

$$\text{e.g. } a_4 = \binom{4}{0} + \binom{4-1}{1} + \binom{4-2}{2} = \binom{4}{0} + \binom{3}{1} + \binom{2}{2} = 1+3+1 = 5$$

This formula is "explicit", but still not so fast to compute
+ still does not give sense of growth.

Also, \exists of this formula doesn't mean there isn't a better one

$$\begin{aligned}\text{e.g. } \# \text{ subsets of } \{1, 2, \dots, n\} &= \sum_{k=0}^n \binom{n}{k} \text{ ("explicit" formula)} \\ &= 2^n \text{ (better formula!)}\end{aligned}$$

③ Explicit formula w/ exponentiation

The recurrence relation $a_n = a_{n-1} + a_{n-2}$ (w/ initial conditions $a_0 = a_1 = 1$)

implies that $a_n = \frac{1}{\sqrt{5}} \left(\underbrace{\left(\frac{1+\sqrt{5}}{2}\right)^{n+1}}_{\varphi} - \underbrace{\left(\frac{1-\sqrt{5}}{2}\right)^{n+1}}_{\bar{\varphi}} \right).$

(You may have seen this e.g. in a linear algebra course.)

We'll explain why this formula holds very soon.

This formula is very explicit, and it does let us see growth rate of a_n ,
but it has its own drawbacks: e.g. why is it even an integer?

④ Asymptotic formula

Can compute $\begin{cases} \varphi \approx 1.618 \dots \text{ ("golden ratio")} \\ \bar{\varphi} \approx -0.618 \dots \end{cases} \Rightarrow |\varphi| > 1$,

which means $a_n \approx \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^{n+1}$, telling us a lot about
its growth rate; e.g., # digits of a_n is

$$\log_{10}(a_n) \approx (n+1) \log_{10}(\varphi) + \log_{10}\left(\frac{1}{\sqrt{5}}\right) \xrightarrow{\text{small}} \approx 0.2$$

⑤ (Ordinary) generating function for a_n

$$A(x) := \underset{\text{DEFN}}{a_0 + a_1 x + a_2 x^2 + a_3 x^3 + \dots} \\ (= 1 + x + 2x^2 + 3x^3 + 5x^4 + \dots)$$

$$= \sum_{n \geq 0} a_n x^n \in \mathbb{C}[[x]],$$

the ring of formal power series in x ,
with coefficients in \mathbb{C} .

(we'll give a formal definition of their
algebraic structure in a short while ...)

Not so clear at first why you'd ever consider $A(x)$,
but we'll see that generating functions are extremely powerful,
e.g.; we can derive everything we saw about (a_n) from $A(x)$.

9/6/19

Claim: $A(x) = \frac{1}{1-x-x^2}$

Pf. Recall
Recurrence $a_n = a_{n-1} + a_{n-2}$ for $n \geq 2$
(and $a_0 = a_1 = 1$)

Multiply by x^n and sum over all $n \geq 2$ to get:

$$\sum_{n \geq 2} a_n x^n = \sum_{n \geq 2} a_{n-1} x^n + \sum_{n \geq 2} a_{n-2} x^n$$

$$A(x) - a_0 x^0 - a_1 x^1 = x \left(\sum_{m \geq 1} a_m x^m \right) + x^2 \left(\sum_{m \geq 2} a_m x^m \right)$$

$$A(x) - 1 - x = x(A(x) - a_0 x^0) + x^2 (A(x))$$

$$(1-x-x^2) A(x) = x + 1 - x = 1$$

$$\Rightarrow A(x) = \frac{1}{1-x-x^2}$$

What good is knowing $A(x) = \frac{1}{1-(x+x^2)}$? Plenty!

Let's extract coefficients of $A(x)$ in various ways...

$$\textcircled{a} \quad A(x) = \frac{1}{1-(x+x^2)} = 1 + (x+x^2) + (x+x^2)^2 + (x+x^2)^3 + \dots$$

$$\begin{aligned} \text{i.e., } \sum_{n \geq 0} a_n x^n &= \sum_{d \geq 0} (x+x^2)^d = \sum_{d \geq 0} \sum_{k=0}^d \binom{d}{k} (x^2)^k x^{d-k} \\ &= \sum_{n \geq 0} x^n \left(\sum_{k=0}^{\lfloor n/2 \rfloor} \binom{n-k}{k} \right) \quad \begin{matrix} \nearrow x^{d+k} \\ \begin{matrix} \nearrow n=d+k \\ d=n-k \end{matrix} \end{matrix} \\ \Rightarrow a_n &= \sum_{k=0}^{\lfloor n/2 \rfloor} \binom{n-k}{k}, \text{ our first explicit formula from before} \end{aligned}$$

$$\textcircled{b} \quad A(x) = \frac{1}{1-x-x^2} = \frac{\frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)}{1 - \frac{1+\sqrt{5}}{2} x} + \frac{\frac{-1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)}{1 - \frac{1-\sqrt{5}}{2} x}$$

How to see this?

$$\text{Recall partial fractions: } \frac{1}{ax^2+bx+c} = \frac{1}{a(x-r_1)(x-r_2)} = \frac{A}{x-r_1} + \frac{B}{x-r_2} \\ = \frac{-A/r_1}{1 - \frac{x}{r_1}} + \frac{-B/r_2}{1 - \frac{x}{r_2}}$$

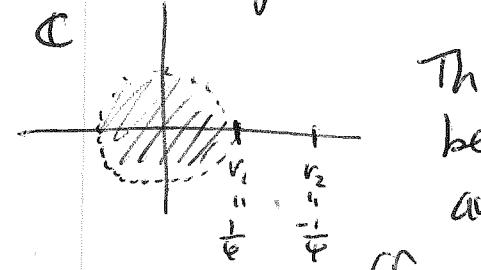
$$\text{Here } r_1 = \left(\frac{1+\sqrt{5}}{2}\right)^{-1}, \quad r_2 = \left(\frac{1-\sqrt{5}}{2}\right)^{-1}$$

$$\rightarrow = \frac{1}{\sqrt{5}} \sum_{n \geq 0} \left(\frac{1+\sqrt{5}}{2}\right)^{n+1} x^n - \frac{1}{\sqrt{5}} \sum_{n \geq 0} \left(\frac{1-\sqrt{5}}{2}\right)^{n+1} x^n$$

$$\rightarrow a_n = \frac{1}{\sqrt{5}} \left(\left(\frac{1+\sqrt{5}}{2}\right)^{n+1} - \left(\frac{1-\sqrt{5}}{2}\right)^{n+1} \right), \text{ our second explicit formula from before...}$$

③ The asymptotic $a_n \approx c \left(\frac{1+\sqrt{5}}{2}\right)^n$ for some constant c

Corresponds to the fact that r_i^{-1} is the reciprocal of the pole of $A(x) = \frac{1}{1-x-x^2} = \frac{1}{(x-r_1)(x-r_2)}$ nearest the origin of \mathbb{C} .



This is just the tip of the rich interplay between thinking of $A(x)$ as a formal power series, and as a function of a complex variable.

(for more see H. Wilf's "generatingfunctionology", linked to on website of the class.)

The fast way to get $A(x) = \frac{1}{1-x-x^2}$ is via Pólya's "picture-writing".

$$\frac{1}{1-(\boxed{} + \boxed{})} = 1 + (\boxed{} + \boxed{})^1 + (\boxed{} + \boxed{})^2 + (\boxed{} + \boxed{})^3 + \dots$$

$\boxed{}$ $\boxed{}$

$n=0 \quad n=1 \quad n=2 \quad \dots$

" " "

$(\boxed{} + \boxed{})(\boxed{} + \boxed{})$ $(\boxed{} + \boxed{})(\boxed{} + \boxed{})(\boxed{} + \boxed{})$

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$\boxed{} + \boxed{} + \boxed{} + \boxed{} + \dots + \boxed{}$

$n=2 \quad n=3 \quad n=3 \quad n=4 \quad \dots$

$\boxed{} \mapsto x^1 \quad \boxed{} \mapsto x^2$

ask $P=x^1$
but $Q=x^2$

$\mathbb{C}[[x]] \ni A(x) = \frac{1}{1-(x+x^2)} \quad (= 1 + (x+x^2)^1 + (x+x^2)^2 + \dots)$

The generating function can often be refined to keep track of additional statistics on our combinatorial objects.

Say we want to compute

$$a_{m,n} = \#\{ \text{tilings of } 2 \times n \text{ rectangle w/ } m \text{ vertical tiles} \}$$

from "picture-writing" we get

$$\left[\frac{1}{1-(P+Q)} \right]_{\substack{P=vx \\ Q=x}} = \sum_{n,m \geq 0} a_{m,n} x^n v^m \in C([v, v])$$

"weight" vertical tile by v , formal parameter

$$\left[\frac{1}{1-(P+Q)} \right]_{\substack{P=vx \\ Q=x}} = \frac{1}{1-vx-x^2}$$

This g.f. is useful for e.g. computing (asymptotically) the expected number of vertical tiles in a random tiling

$$\begin{aligned} \sum_{n \geq 0} \left(\sum_{m \geq 0} a_{m,n} \cdot m \right) x^n &= \left[\frac{\partial}{\partial v} \sum_{n,m \geq 0} a_{m,n} x^n v^m \right]_{v=1} \\ &= \left[\frac{\partial}{\partial v} \frac{1}{1-vx-x^2} \right]_{v=1} \\ &= \left[\frac{x}{(1-vx-x^2)^2} \right]_{v=1} = \frac{x}{(1-x-x^2)^2} \end{aligned}$$

Via partial fractions $\frac{x}{(1-x-x^2)^2} = \frac{A_1 x + B_1}{(x-r_1)^2} + \frac{A_2 x + B_2}{(x-r_2)^2} + \frac{C}{(x-r_1)} + \frac{D}{(x-r_2)}$

(can use above formula to show $\sum_{m \geq 0} a_{m,n} m \approx \frac{n}{5} \left(\frac{1+\sqrt{5}}{2}\right)^{n+1}$)

Recall $a_n \approx \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^{n+1} \Rightarrow \sum_{m \geq 0} a_{m,n} m \approx \frac{n}{\sqrt{5}} \cdot a_n$

Thus, the expected # of vertical tiles is $\approx \frac{n}{\sqrt{5}}$, i.e., out of the n tiles, about $\frac{1}{\sqrt{5}}$ of them will be vertical.

9/9/2019

The ring of formal power series $R[[x]]$

(Where $R = \mathbb{C}$ or \mathbb{R} or \mathbb{Q} or $\mathbb{C}[v]$ or any commutative ring w/ 1)
or \mathbb{H}_q

DEF'N $R[[x]] := \left\{ a_0 + a_1 x + a_2 x^2 + \dots = \sum_{n=0}^{\infty} a_n x^n \text{ w/ } (a_0, a_1, a_2, \dots) \in R \right\}$
is a commutative ring w/ coefficientwise
addition: $A(x) + B(x) = \sum_{n=0}^{\infty} (a_n + b_n) x^n$,
and multiplication via convolution:

$$(C(x)) := A(x)B(x) = \sum_{n=0}^{\infty} c_n x^n \text{ where } c_n = \sum_{i=0}^n a_i b_{n-i}$$
$$= a_0 b_0 + (a_0 b_1 + a_1 b_0) x + (a_0 b_2 + a_1 b_1 + a_2 b_0) x^2 + \dots$$

So its zero is $0 = 0 + 0x + 0x^2 + \dots$

and its one is $1 = 1 + 0x + 0x^2 + 0x^3 + \dots$

Prop. $A(x) = \sum_{n=0}^{\infty} a_n \in R[[x]]$ is a unit (i.e., $\exists B(x)$ w/ $1 = A(x)B(x)$)
 $\iff a_0$ is a unit of R (i.e., $\exists b_0 \in R$ w/ $1 = a_0 b_0$).

E.g. By this criterion, $(1-x-x^2) \in \mathbb{C}[[x]]$ is a unit, so
 $\exists A(x)$ w/ $A(x)(1-x-x^2)=1$, i.e. $A(x) = \frac{1}{1-x-x^2}$ ($= 1+x+2x^2+3x^3+5x^4+\dots$)

Proof:

$$1 = A(x)B(x) = a_0 b_0 + (a_0 b_1 + a_1 b_0) x + (a_0 b_2 + a_1 b_1 + a_2 b_0) x^2 + \dots$$
$$\iff a_0 b_0 = 1 \quad (\text{so } a_0 \text{ needs to be a unit})$$

(i.e., $b_0 = a_0^{-1}$ in R)

and then

$$a_0 b_1 + a_1 b_0 = 0 \text{ means } b_1 = \frac{-a_1 b_0}{a_0} \quad \begin{matrix} \text{allowed since } b_0 = a_0^{-1} \\ \text{already defined} \end{matrix}$$

$$a_0 b_2 + a_1 b_1 + a_2 b_0 = 0 \text{ means } b_2 = \frac{-(a_1 b_1 + a_2 b_0)}{a_0}$$

... (can recursively compute all b_i in unique way) 

DEF'N A sequence $A_1(x), A_2(x), \dots$ in $\mathbb{R}[[x]]$ converges (i.e., $\exists A(x)$ w/ $A(x) = \lim_{y \rightarrow \infty} A_y(x)$) if $\forall n \geq 0$, the coefficient of x^n in $A_j(x)$ stabilizes for $j \gg 0$:

denote this by $[x^n]A_j(x)$ i.e., $\forall n \geq 0, \exists N > 0$ and $a_n \in \mathbb{R}$ s.t. $[x^n]A_j(x) = a_n \quad \forall j \geq N.$

$$A(x) = \frac{1}{1-(x+x^2)} = \underbrace{1}_{A_0(x)} + \underbrace{(x+x^2)}_{A_1(x)} + \underbrace{(x+x^2)^2}_{A_2(x)} + \underbrace{(x+x^2)^3}_{A_3(x)} + \dots$$

Converges in $\mathbb{C}[[x]]$, e.g. $[x^3]A(x) = [x^3]A_3(x) = [x^3]A_4(x) = \dots = a_3 = 3,$

$e^{x+1} := 1 + \frac{(x+1)}{1!} + \frac{(x+1)^2}{2!} + \frac{(x+1)^3}{3!}$ does not converge ($x+1 \notin \mathbb{R}(x)$)
 but $e^x := 1 + \frac{x}{1!} + \frac{x^2}{2!} + \dots$ does converge.

Alternatively, $\{A_j(x)\}_{j=0,1,\dots}$ converges in $\mathbb{R}[[x]]$,

if $\lim_{j \rightarrow \infty} \min \deg(A_j(x) - A_{j-1}(x)) = \infty$

where DEF'N $\text{mdeg } A(x) :=$ $\begin{cases} \text{smallest } d \text{ w/ } ad \neq 0 \\ \sum_{n \geq 0} a_n x^n \end{cases}$ ($= \infty$ if no such d)

e.g. $A(x) = \frac{1}{1-x-x^2}$, then $A_j(x) - A_{j-1}(x) = (x+x^2)^j$
 and $\min \deg((x+x^2)^j) = j \rightarrow \infty$ as $j \rightarrow \infty$.

Cor $\sum_{j=0}^{\infty} B_j(x) = B_0(x) + B_1(x) + B_2(x) + \dots$ converges in $R[[x]]$

$\underbrace{B_0(x)}_{\text{if } \lim_{n \rightarrow \infty} A_n(x)}$ $\underbrace{B_1(x)}_{(w/B_j = A_j - A_{j-1})}$ $\underbrace{B_2(x)}_{A_2(x)}$

$\Leftrightarrow \min \deg B_j(x) \rightarrow \infty$ as $j \rightarrow \infty$.

Cor Infinite products of the form

$$\prod_{j=1}^{\infty} (1 + B_j(x)) \text{ w/ } \min \deg B_j \geq 1 \quad (\forall j)$$

Converges in $R[[x]] \Leftrightarrow \lim_{j \rightarrow \infty} \min \deg B_j = \infty$

$$(\text{=} \lim_{n \rightarrow \infty} A_n(x) \text{ where } A_0 = 1, A_1 = (1 + B_1(x)), \\ A_2 = (1 + B_1(x))(1 + B_2(x)), \dots)$$

E.g. $\prod_{n=1}^{\infty} (1 + \frac{1}{2^n} x) \text{ does not converge in } R[[x]]$

(even if it does make sense analytically to think of $A(x)$ as a function of $x \in \mathbb{C}$ for at least $|x|$ small enough)

E.g. $\prod_{n=1}^{\infty} (1 + x^n)$ converges in $R[[x]]$

$$(1+x)(1+x^2)(1+x^3)(1+x^4)(1+x^5)\dots$$

$$\sum_{n=0}^{\infty} a_n x^n$$

Q: What are these coefficients a_n ?

A: We'll see next time.

11/16/19 Partition generating functions

DEF'N A (number) partition $\lambda = (\lambda_1, \lambda_2, \lambda_3, \dots)$ of n is a weakly decreasing sequence of nonnegative integers eventually $\lambda_1 \geq \lambda_2 \geq \lambda_3 \geq \dots \geq 0$ with sum $\lambda_1 + \lambda_2 + \dots = 0$, (i.e., $\lambda_i \in \mathbb{N}$ $\forall i = 1, 2, \dots$)

We write $\lambda \vdash n$ (λ $\backslash\!\!\!/\!\!\!\dashv n$) and $|\lambda| = n$, "size" of λ .

e.g. $\lambda = (5, 5, 3, 1, 0, 0, \dots) = (5, 5, 3, 1, 0) = (5, 5, 3, 1) \vdash 14 = 5+5+3+1$,

its length $l(\lambda) := \{i : \lambda_i > 0\} = \# \text{ of nonzero parts } \lambda_i$

its Young diagram is a left + top justified array of boxes with λ_i boxes in the i^{th} row from the top:

e.g. $\lambda = (5, 3, 3, 1) \longleftrightarrow$

Let $p(n) := \# \text{ of partitions } \lambda \vdash n$

	p(n)	n
	1	7
	7	5
	5	4
	3	3
	2	2
	1	1
\emptyset (empty partition, unique partition of 0)	1	0

$\mathcal{Y} = \text{Young's lattice}$, the poset of all partitions ordered by containment of Young diagrams

$$\sum_{n \geq 0} p(n) q^n = \sum_{\substack{\text{all partitions} \\ X}} q^{|X|} = \left[\begin{array}{ll} ((1+A_1+A_1^2+A_1^3+\dots) & ((1+A_2+A_2^2+A_2^3+\dots)) \\ (q^0+q^1+q^2+q^3+\dots) & (q+q^2+q^4+q^6+\dots) \\ ((1+D)+\boxed{1}+\boxed{1}+\dots) & (1+\boxed{\square}+\boxed{\square}+\boxed{\square}+\dots) \\ ((1+A_3+A_3^2+\dots) & (1+A_4+\dots) \\ (1+q^3+q^6+\dots) & (1+q^4+q^8+\dots) \\ (0+\boxed{1}+\boxed{1}+\dots) & (1+\boxed{\square\square}+\boxed{\square\square}+\dots) \end{array} \right]$$

\vdots

$$= \left[\begin{array}{l} + \dots + \begin{matrix} A_4 \\ A_3 \\ A_2 \\ A_2 \\ A_2 \\ A_1 \end{matrix} = A_1^2 A_2^4 A_3^0 A_4^4 \end{array} \right]$$

$A_i \mapsto q^i$

$$= ((1+q+q^2+q^3+\dots)(1+q^2+(q^2)^2+\dots)(1+q^3+(q^3)^2+\dots)\dots)$$

$\xrightarrow{B_1} \quad \xrightarrow{B_2} \quad \xrightarrow{B_3}$

$$= \frac{1}{1-q} \cdot \frac{1}{1-q^2} \cdot \frac{1}{1-q^3} \cdots \text{ a convergent product!}$$

$$= \prod_{n=1}^{\infty} \frac{1}{(1-q^n)} \quad \leftarrow \text{g.f. for all partitions as a product formula!}$$

Cultural asides:

Thm (Euler's pentagonal # thm)

$$\prod_{n=1}^{\infty} (1-q^n) = 1 + \sum_{k=1}^{\infty} (-1)^k (x^{k(3k+1)/2} + x^{k(3k-1)/2})$$

"pentagonal #'s"

Pf: Look up. Either Euler's original algebraic proof via g.f.'s, or very nice bijection pf due to Franklin.

Together with above g.f. for partitions get $\sum p(n) q^n = (\prod (1-q^n))^{\prime \prime}$

$$\Rightarrow \text{Recurrence for } p(n) = \sum_{\substack{k=3+0, k \neq 0}}^{\infty} (-1)^k p(n-k(3k-1)/2) \quad \leftarrow P(m)=0 \text{ if } m < 0$$

$$= p(n-1) + p(n-2) - p(n-5) - p(n-7) + p(n-12) + p(n-13) + \dots$$

Thm (Hardy + Ramanujan) $p(n) \approx \frac{1}{4n\sqrt{3}} e^{\pi \sqrt{\frac{2n}{3}}}$

Pf: "circle method", careful analysis of singularities of $\prod \frac{1}{(1-q^n)}$, residues, etc.

G.f.'s for restricted partitions:

Let $q(n) := \#$ of partitions of n into distinct parts.

n	$q(n)$	
0	1	\emptyset
1	1	\square
2	1	$\square\square$
3	2	$\square\square\square$ $\square\square\square$
4	2	$\square\square\square\square$ $\square\square\square\square$ $\square\square\square\square$
5	3	$\square\square\square\square\square$ $\square\square\square\square\square$ $\square\square\square\square\square$ $\square\square\square\square\square$

$$Q(q) := \sum_{n \geq 0} q(n) q^n = (1+q)(1+q^2)(1+q^3)(1+q^4) \dots$$

$$= \prod_{i \geq 1} (1+q^i) \quad \text{converges in } \mathbb{C}(q)$$

$\text{et } P_{\text{odd}}^{(m)} = \# \text{ partitions of } n \text{ into odd parts}$

n	$P_{\text{odd}}(n)$	
0	1	\emptyset
1	1	\square
2	1	\square
3	2	$\square\square\square$
4	2	$\square\square\square$
5	3	$\square\square\square\square\square$

Looks like $q(n) = P_{\text{odd}}(n)$, but how to show?

$$\text{G.f.'s! } P_{\text{odd}}(q) = \sum_{n \geq 0} P_{\text{odd}}(n) q^n = (1+q+q^2+\dots)(1+q^3+q^6+\dots)$$

$$(1+q^5+q^{10}+\dots) \quad \text{converges in } \mathbb{C}(q) \quad \frac{1}{(1-q)(1-q^3)(1-q^5)\dots} = \frac{1}{\prod_{i \geq 0} (1-q^{2i+1})} \left(\prod_{i=0}^{\infty} \prod_{j=1}^{2i+1} (1+q^j) \right)$$

How to show $Q(q) = \text{Podd}(q)$?

$$\begin{aligned} \text{Well, } Q(q) &= (1+q)(1+q^2)(1+q^3) \dots \\ &= \frac{(1-q^2)}{(1-q)} \cdot \frac{1+(q^2)^2}{1-q^2} \cdot \frac{1-(q^3)^2}{1-q^3} \dots \\ &= \frac{(1-q^3)(1+q^4)}{(1-q)(1+q^2)(1-q^3)(1+q^4)} \dots \\ &= \frac{1}{(1-q)(1-q^3)(1-q^5)} \dots = \text{Podd}(q)! \end{aligned}$$

Was that manipulation allowed? Yes! Let's think of it slightly differently...

$$\text{Let } R(q) := (1-q)(1-q^3)(1-q^5) \dots = \frac{1}{\text{Podd}(q)} \in \mathbb{C}[[q]],$$

Want to show $1 \neq Q(q)R(q)$ in $\mathbb{Q}[[q]]$:

$$\begin{aligned} (1+0 \cdot q + 0 \cdot q^2 \dots) &\quad \left(\underbrace{(1+q)(1+q^2)(1+q^3) \dots}_{\text{converges}} \right) \left(\underbrace{(1-q)(1-q^3)(1-q^5) \dots}_{\text{converges}} \right) \\ &= \underbrace{(1+q)(1-q)}_{(1-q^2)} \cdot \underbrace{(1+q^2)(1-q^3) \dots}_{(1-q^4)} \cdot \underbrace{(1+q^4)(1-q^5) \dots}_{(1-q^6)} \dots \\ &= (1-q^4) \cdot \underbrace{(1+q^3)(1+q^4) \dots}_{(1+q^5)} \cdot \underbrace{(1-q^3)(1-q^5) \dots}_{(1-q^7)} \dots \\ &= \underbrace{(1-q^4)(1-q^6)}_{(1-q^8)} \cdot \underbrace{(1+q^4)(1+q^5) \dots}_{(1+q^7)} \cdot \underbrace{(1-q^5)(1-q^7) \dots}_{(1-q^9)} \dots \\ &= (1-q^8)(1-q^6) \cdot \underbrace{(1+q^5)(1+q^6) \dots}_{(1+q^7)} \cdot \underbrace{(1-q^5)(1-q^7) \dots}_{(1-q^9)} \dots \\ &= \dots \text{ etc.} \quad = 1 + 0 \cdot q + 0 \cdot q^2 + 0 \cdot q^3 + 0 \cdot q^4 \dots \end{aligned}$$

↑ Bijective proofs of $q(n) = \text{Podd}(n)$ as well

(See Stanley Prop. 1.8.5)
for two such proofs!

↑ hinted at
by (this)
manipulation

$$\begin{aligned} \text{e.g. } \lambda &= (9^5, 5^2, 3^2, 1^3) = (q^{2+2^2}, q^{2+2^2}, q^2, q^{2+2^1}) \text{ & } \mu = (9 \cdot 2^0, 9 \cdot 2^2, 5 \cdot 2^2, \\ &\quad \text{Podd}(n) \quad (9, 9, 9, 9, 5, 5, 3, 3, 1, 1, 1) \quad \text{etc.} \rightarrow 5 \cdot 2^3, 5 \cdot 2^3, 3 \cdot 2^1, 1 \cdot 2^0, 1^1) \\ &\quad \text{etc.} \rightarrow (9, 36, 20, 40, 6, 1, 2) \end{aligned}$$

Some more about formal power series + ordinary generating fun's.

Let's define some specific elements of $\mathbb{C}[[x]]$:

DEFN $e^x := \sum_{n \geq 0} \frac{x^n}{n!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots$

$$\log(1+x) := \sum_{n \geq 1} (-1)^{n-1} \frac{x^n}{n} = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots$$

$$\forall \lambda \in \mathbb{C}, (1+x)^\lambda := \sum_{k \geq 0} \binom{\lambda}{k} x^k,$$

where $\binom{\lambda}{k} := \frac{\lambda(\lambda-1)(\lambda-2)\dots(\lambda-(k-1))}{k!} \in \mathbb{C}$

(Just like $\binom{n}{k} = \frac{n!}{k!(n-k)!} = \frac{n(n-1)\dots(n-(k-1))}{k!}$ for $n \in \mathbb{N}$)

These satisfy the properties you would expect..

e.g. ① ~~$(1+x)^\lambda (1+x)^M = (1+x)^{\lambda+M} \in \mathbb{C}[[x]]$~~

② $e^x e^y = e^{x+y} \in \mathbb{C}[[x,y]]$, ...

③ $e^{\log(1+x)} = 1 + x$, etc...

defined to be $= 1 + \log(1+x) + \frac{(\log(1+x))^2}{2!} + \dots$

$$= 1 + \left(x - \frac{x^2}{2} + \frac{x^3}{3} - \dots\right) + \frac{\left(x - \frac{x^2}{2} + \frac{x^3}{3} - \dots\right)^2}{2!} + \dots$$

why does this even converge in $\mathbb{C}[[x]]$?

General Prop. If $A(x) = \sum_{n \geq 0} a_n x^n$, $B(x) = \sum_{n \geq 0} b_n x^n$, and $b_0 = 0$,

then $A(B(x)) := \sum_{n \geq 0} a_n (B(x))^n$ converges in $\mathbb{C}[[x]]$.

Note: $\log(1+x) = x - \frac{x^2}{2} + \dots$ has 0 constant term.

How to justify ①, ②, ③, etc...? Could do very tedious manipulation of coefficients, but instead, track from analysis^{complex}, some $R > 0$

Thm If $f(z) = \sum a_n z^n$ is analytic for $|z| < R$ and
 ↗ if vanishes on $|z| < R$, then $f(z) \equiv 0$, i.e., $a_0 = a_1 = a_2 = \dots = 0$.
 (Standard fact from complex analysis, true under much weaker assumptions)

④ for $n \in \mathbb{N}$, $(1+x)^n = \sum_{k=0}^n \binom{n}{k} x^k = \sum_{k=0}^n \binom{n}{k} x^k$

but also $\frac{1}{(1-x)^n} = (1+(-x))^n = \sum_{k=0}^n \frac{(-n)(-n-1)(-n-2)\dots(-n-(k-1))}{k!} (-1)^k x^k$
 "
 $(1+x+x^2+\dots)(1+x+x^2+\dots)\dots(1+kx+x^2+\dots)$ $\overset{\text{as above}}{=} \sum_{k=0}^n \frac{n(n+1)(n+2)\dots(n+k-1)}{k!} x^k$
 onion $\overset{\text{same}}{\Rightarrow}$ $\overset{\text{n= # parentheses}}{\Rightarrow}$ $\overset{\text{# flavors}}{\Rightarrow}$ $\overset{\text{of bagels}}{\Rightarrow} \sum_{k=0}^n \binom{n+k-1}{k} x^k$
 $\Rightarrow \# \text{stars} / \# \text{bars} = \# \text{flavors of bagels}$

Pf that $\binom{n}{k} = \binom{n+k-1}{k}$:

"Stars and bars"

stars indicate how many times

each element is chosen, bar separate the bins that represent elements.

⑤ $\frac{1}{1-4x} = \sum_{k=0}^{\infty} \binom{-1}{k} (-4x)^k = \sum_{k=0}^{\infty} \underbrace{\binom{1+k-1}{k}}_{\binom{k}{k}=1} 4^k x^k = \sum_{k=0}^{\infty} 4^k x^k$

but also $\frac{1}{(1-4x)^2} = \sum_{k=0}^{\infty} \underbrace{\binom{2+k-1}{k}}_{\binom{k+1}{k}=k+1} 4^k x^k = \sum_{k=0}^{\infty} (k+1) 4^k x^k$

$\frac{1}{(1-4x)^3} = \sum_{k=0}^{\infty} 4^k x^k \cdot \binom{k+2}{2}, \quad \frac{1}{(1-4x)^4} = \dots$ etc.

Useful for extracting coefficients after partial fraction expansions.

$$\begin{aligned}
 ⑥ \frac{1}{\sqrt{1-4x}} &= (1-4x)^{-\frac{1}{2}} = \sum_{k \geq 0} \binom{-\frac{1}{2}}{k} (-4)^k x^k \\
 &= \sum_{k \geq 0} \left(\frac{\frac{-1}{2}(\frac{-3}{2})(\frac{-5}{2}) \cdots \frac{-(2k-1)}{2}}{k!} (-4)^k \right) x^k \\
 &= \sum_{k \geq 0} \left(\frac{4^k (1)(3)(5) \cdots (2k-1)}{2^k k!} \right) x^k \\
 &= \sum_{k \geq 0} 2^k \frac{(2k)!}{k! k!} \frac{(1)(3)(5) \cdots (2k-1)}{k!} x^k = \sum_{k \geq 0} \frac{(2k)!}{k! k!} x^k
 \end{aligned}$$

interesting... $\rightarrow \sum_{k \geq 0} \binom{2k}{k} x^k$

Another tool from calculus for $R[[x]]$...

DEF'N for $A(x) = \sum_{n \geq 0} a_n x^n \in R[[x]]$, we define
 the formal derivative $A'(x) := \sum_{n \geq 1} \underbrace{n a_n}_{\substack{a_1 + a_2 + \cdots + a_n \\ n \text{ times}}} x^{n-1} \in R[[x]]$.

The derivative satisfies usual rules from calculus;

$$(A(x) + B(x))' = A'(x) + B'(x)$$

$$(AB)' = A'B + B'A$$

$$\left(\frac{1}{A}\right)' = -\frac{A'}{A^2}$$

$$(A(B(x)))' = A'(B(x)) \cdot B'(x)$$

- etc...

Quick review of binomials and multinomials

Binomial coefficient $\binom{n}{k} = \frac{n!}{k!(n-k)!}$ has several (easy) interpretations

($\binom{n}{x, n-x}$)
(in multinomial notation)

= # words with k 1's

e.g. $\binom{9}{2} = 6 \rightarrow$

$n-k$ 2's

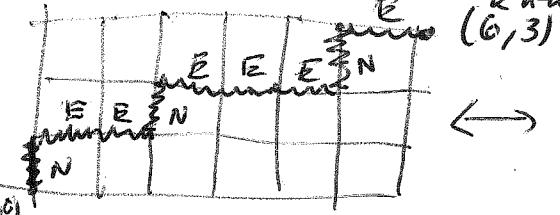
{1122, 1212, 1221, 3
2121, 2112, 2211}

i.e., rearrangements of $\underbrace{111\dots 1}_k \underbrace{222\dots 2}_{n-k}$

= # lattice paths in \mathbb{Z}^2 taking E or N steps,
from $(0,0)$ to $(k, n-k)$

e.g.

$n=9$
 $k=6$
 $n-k=3$

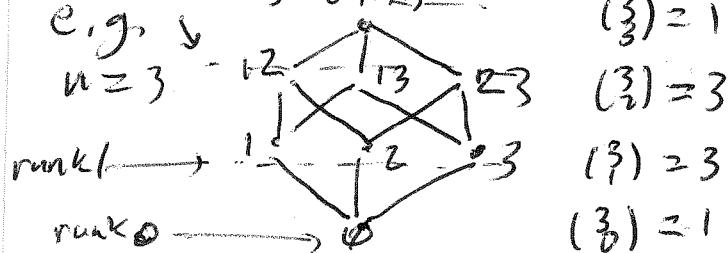


\leftarrow

$\binom{9}{6,3}$
 E^6, N^3

= size of the k^{th} rank in B_n , where B_n
is the Boolean lattice of subsets of $\{1, 2, \dots, n\}$
partially ordered by containment

rank 3: $\binom{3}{1,2,3}$



$\binom{3}{3} = 1$

$\binom{3}{2} = 3$

$\binom{3}{1} = 3$

$\binom{3}{0} = 1$

Of course, also have binomial theorem: $(x+y)^n = \sum_{k=0}^n \binom{n}{k} xy^{n-k}$

Multinomials:

How many rearrangements of BANANA'S?

i.e., of 3 A's, 1 B, 2 N's, 1 S? (equiv. of $\overset{1}{A}\overset{2}{A}\overset{3}{A}\overset{4}{B}\overset{5}{N}\overset{6}{N}\overset{7}{S}$)

The transitive action of the symmetric group G_7 on the rearrangements

e.g., permutation $\sigma = (1\ 2\ 3\ 4\ 5\ 6\ 7)$ sends AAA \bar{B} NNNN \mapsto AABASNN

The stabilizer of this action is $G_3 \times G_4 \times G_2 \times G_1 \subseteq G_7 \rightarrow (7, 1, 2, 1)$

So by orbit-stab. thm, ~~$|G_7|$~~ = $\frac{|G_7|}{|\text{# arrangements}|} = \frac{7!}{(G_3 \times G_4 \times G_2 \times G_1)} = \frac{7!}{3!(1)(2)1!}$

binomial coefficient $\binom{n}{k_1, k_2, \dots, k_m} = \frac{n!}{k_1! k_2! \dots k_m!} \text{ for } k_1 + \dots + k_m = n$

= # words w/ $\begin{matrix} k_1 1's \\ k_2 2's \\ \vdots \\ k_m m's \end{matrix}$, i.e. rearrangements of $\underbrace{1 \dots 1}_{k_1} \underbrace{2 \dots 2}_{k_2} \dots \underbrace{m \dots m}_{k_m}$

= # lattice paths in \mathbb{Z}^m taking steps e_1, e_2, \dots, e_m (standard basis) from $0 = (0, 0, \dots, 0)$ to (k_1, k_2, \dots, k_m)

(same correspondence between words + paths as binomials)

\Rightarrow # chains (or flags) in B_n of subsets

$$\phi: S_0 \subseteq S_{k_1} \subset S_{k_1+k_2} \subset \dots \subset S_{n+k_2+\dots+k_{m-1}} \subset S_n = \{1, 2, \dots, n\}$$

passing through ranks $0, k_1, k_1+k_2, \dots, k_1+k_2+\dots+k_{m-1}, n$.

e.g. for $(3, 1, 2, 1)$ have 2131314 $\leftrightarrow \phi \subset \overbrace{S_3}^{\{1, 2, 3\}} \subset \overbrace{S_4}^{\{1, 2, 3, 4\}} \subset \overbrace{S_6}^{\{1, 2, 3, 4, 5, 6\}} \subset [7]$.

Note: $\binom{n}{k_1, k_2, \dots, k_m} = \binom{n}{k_1} \binom{n-k_1}{k_2} \binom{n-(k_1+k_2)}{k_3} \dots \binom{n-(k_1+\dots+k_{m-2})}{k_{m-1}} \binom{k_m}{k_m}$

binomial theorem: $(x_1 + x_2 + \dots + x_m)^n = \sum_{\substack{(k_1, k_2, \dots, k_m) \\ k_1 + \dots + k_m = n}} \binom{n}{k_1, k_2, \dots, k_m} x_1^{k_1} x_2^{k_2} \dots x_m^{k_m}$